### **Chapter 17**

### Methodology Logical Database Design for the Relational Model

### Logical Database Design for the Relational Model

- Step 2 Using a Data Model (e.g., Relational Data Model), Build logical data model and validate it
  - Step 2.1 Derive relations for logical data model
  - Step 2.2 Validate relations using normalization
  - Step 2.3 Check integrity constraints

 Create relations for logical data model to represent entities, relationships, and attributes that have been identified

- In this step, we derive relations for the logical data model to represent the entities, relationships, and attributes
- We describe the composition of each relation using a Database Definition Language (DBDL) for relational databases
- Using the DBDL, we first specify the name of the relation followed by a list of the relation's simple attributes enclosed in brackets
- We then identify the primary key and any alternate and/or foreign key(s) of the relation
- Following the identification of a foreign key, the relation containing the referenced primary key is given
- Any derived attributes are also listed together with how each one is calculated

- The relationship that an entity has with another entity is represented by the primary key/foreign key mechanism
- In deciding where to post (or place) the foreign key attribute(s), we must first identify the 'parent' and 'child' entities involved in the relationship
- The parent entity refers to the entity that posts a copy of its primary key into the relation that represents the child entity, to act as the foreign key

- We describe how relations are derived for the following structures that may occur in a conceptual data model:
- (1) strong entity types;
- (2) weak entity types;
- (3) one-to-many (1:\*) binary relationship types;
- (4) one-to-one (1:1) binary relationship types;
- (5) one-to-one (1:1) recursive relationship types;
- (6) superclass/subclass relationship types;
- (7) many-to-many (\*:\*) binary relationship types;
- (8) multi-valued attributes.

#### **Conceptual Data Model of all Users' Views**



- (1) Strong entity types
  - For each strong entity in data model,
    - Create **relation** that includes all simple attributes of entity
    - For **composite attributes**, include **only constituent simple** attributes: fName and IName in the relation

**Staff** (staffNo, fName, IName, position, gender, DOB) **Primary Key** staffNo

### (2) Weak entity types

- For each weak entity in data model,
  - Create relation that includes all simple attributes of entity
  - Primary key of weak entity is partially or fully derived from each owner entity

**Preference (**prefType, maxRent) **Primary Key** None (at present)

- (3) One-to-many (1:\*) binary relationship types
  - For each 1:\* binary relationship,
    - Entity on 'one side' of relationship designated as parent entity
    - Entity on 'many side' designated as child entity
  - Represent relationship by posting copy of primary key attribute(s) of parent entity into relation representing child entity
    - Acts as foreign key

• (3) One-to-many (1:\*) binary relationship types

Post staffNo into Client to model 1:\* Registers relationship

Staff (staffNo, fName, IName, position, sex, DOB) Primary Key staffNo

Client (clientNo, fName, IName, telNo, staffNo) Primary Key clientNo Foreign Key staffNo references Staff(staffNo)

- (4) One-to-one (1:1) binary relationship types
  - Cardinality cannot be used to identify parent and child entities
  - **Participation constraints** are used instead to decide:
  - Options:

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- 1) Combine entities involved into one relation
- 2) Create two relations and post copy of primary key from one relation to other
- Consider the following:
  - *(a) mandatory* participation (1:1 vs 1:1)on *both* sides of 1:1 relationship
  - (b) mandatory participation (1:1 vs 0:1) on one side of 1:1 relationship
  - *(c) optional* participation (0:1 vs 0:1) on *both* sides of 1:1 relationship

- (a) *Mandatory* participation (1:1 vs 1:1) on *both* sides of 1:1 relationship
  - Combine entities involved into 1 relation
  - Choose one primary key of original entities to be primary key of new relation
  - Other primary key (if one exists) used as alternate key

- (a) *Mandatory* participation (1:1 vs 1:1) on *both* sides of 1:1 relationship example:
- The Client States Preference relationship is an example of a 1:1 relationship with mandatory participation on both sides

**Client** (clientNo, fName, IName, telNo, prefType, maxRent, staffNo)

- Primary Key clientNo
- Foreign Key staffNo references Staff(staffNo)

- (b) Mandatory participation (1:1 vs 0:1) on one side of a 1:1 relationship
  - Identify parent and child entities using participation constraints (assuming not every client specifies preferences)
  - Entity with optional participation side in relationship designated as child
  - Entity with mandatory participation designated as parent
  - Place copy of primary key of parent in relation representing child

For 1:1 relationship with mandatory participation on Client side, post clientNo into Preference to model States relationship

Client (clientNo, fName, IName, telNo, staffNo)

Primary Key clientNo

Foreign Key staffNo references Staff(staffNo)

Preference (clientNo, prefType, maxRent) Primary Key clientNo Foreign Key clientNo references Client(clientNo)

- (c) Optional participation (0:1 vs 0:1) on both sides of a 1:1 relationship
  - Designation of parent and child entities arbitrary
  - 'Staff Uses Car' Example

- (5) one-to-one (1:1) recursive relationship types;
- For a 1:1 recursive relationship, follow the rules for participation as described previously for a 1:1 relationship
- However, in this special case of a 1:1 relationship, the entity on **both sides** of the relationship is the **same**
- For a 1:1 recursive relationship with mandatory participation on both sides, represent the recursive relationship as a single relation with two copies of the primary key
- One copy of the primary key represents a foreign key and should be renamed to indicate the relationship it represents e.g. manager\_id

- (5) one-to-one (1:1) recursive relationship types;
- For a 1:1 recursive relationship with mandatory participation on only one side, we have the option to create a single relation with two copies of the primary key as described previously,

OR

- to create a new relation to represent the relationship
- The new relation would have only two attributes, both copies of the primary key
- For a 1:1 recursive relationship with optional participation on both sides, again create a new relation as described earlier

• (6) superclass/subclass relationship types;

As described earlier in case of EERM

- (7) Many-to-many (\*:\*) binary relationship types
  - For each \*:\* binary relationship create a relation to represent relationship (associative entity) and include any attributes that are part of relationship
  - Post copy of primary key attribute(s) of entities that participate in relationship into new relation - act as foreign keys
  - Foreign keys also form primary key of new relation
    - Possibly in combination with other attributes of relationship

#### (7) Many-to-many (\*:\*) binary relationship types

Client (clientNo, fName, IName, telNo, prefType, maxRent, staffNo) Primary Key clientNo

Foreign Key staffNo references Staff(staffNo)

PropertyForRent (propertyNo, street, city, postcode, type, rooms, rent)

Primary Key propertyNo

Viewing (clientNo, propertyNo, dateView, comment) Primary Key clientNo, propertyNo Foreign Key clientNo references Client(clientNo) Foreign Key propertyNo references PropertyForRent(propertyNo)

#### • (8) Multi-valued attributes

- Create new relation to represent multi-valued attribute
- Include primary key of main entity in new relation acts as foreign key
- Unless the multi-valued attribute is itself an alternate key of the entity, the primary key of the new relation is the combination of the multivalued attribute and the primary key of the main entity

### (8) Multi-valued attributes

Post branchNo into Telephone

Branch (branchNo, street, city, postcode) Primary Key branchNo Telephone (telNo, branchNo) Primary Key telNo Foreign Key branchNo references Branch(branchNo)

## Summary of how to map entities and relationships to relations

Entity/Relationship	Mapping
Strong entity	Create relation that includes all simple attributes.
Weak entity	Create relation that includes all simple attributes (primary key still has to be identified after the relationship with each owner entity has been mapped).
1:* binary relationship	Post primary key of entity on 'one' side to act as foreign key in relation representing entity on 'many' side. Any attributes of relationship are also posted to 'many' side.
1:1 binary relationship:	
<ul><li>(a) Mandatory participation on both sides</li><li>(b) Mandatory participation on one side</li></ul>	Combine entities into one relation. Post primary key of entity on 'optional' side to act as foreign key in relation representing entity on 'mandatory' side.
(c) Optional participation on both sides	Arbitrary without further information.
Superclass/subclass relationship	See Table 16.1.
*:* binary relationship, complex relationship	Create a relation to represent the relationship and include any attributes of the relationship. Post a copy of the primary keys from each of the owner entities into the new relation to act as foreign keys.
Multi-valued attribute	Create a relation to represent the multi-valued attribute and post a copy of the primary key of the owner entity into the new relation to act as a foreign key.

- Document relations and foreign key attributes
  - At the end of Step 2.1, document the composition of the relations derived for the logical data model using the DBDL
  - Double check all primary and foreign keys again

#### **Conceptual Data Model of all Users' Views**



#### Creating tables to represent specialization/generalization



#### Relations for the *DreamHome* example (Relational Schema)

Manager (staffNo, mgrStartDate, Bonus)	Branch (branchNo,street,city,postcode,mgrNo)
Primary Key staffNo	Primary Key branchNo
Foreign Key staffNo references Staff(staffNo)	Foreign Key mgrNo references Manger(staffNo)
Telephone (telNO, branchNo)	Staff (staffNo, fName, IName, position, gender, DOB, supervisorNo, branchNo)
Primary Key telNO	Primary Key staffNo
Foreign Key branchNo references Branch(branchNo)	Foreign Key supervisorNo references Staff(staffNo)
	Foreign Key branchNo references Branch(branchNo)
Car (carNo, model, companyName ,staffNo)	Client (clientNo, fName, IName, telNo, prefType, maxRent, housNo, streetNo,
Primary Key carNo	zipCode, city, country, staffNo)
Foreign Key staffNo references Staff(staffNo)	Primary Key clientNo
	Foreign Key staffNo references Staff(staffNo)
Viewing (clientNo, propertyNo, dateView, comment)	Registration (clientNo, branchNo, staffNo, dateRegistered)
Primary Key clientNo, propertyNo	Primary Key clientNo, branchNo, staffNo, dateRegistered
Foreign Key clientNo references Client(clientNo)	Foreign Key clientNo references Client(clientNo)
Foreign Key propertyNo references PropertyForRent(propertyNo)	Foreign Key branchNo references Branch(branchNo)
	Foreign Key staffNo references Staff(staffNo)
PropertyForRent (propertyNo, street, city, postcode, rent, room, type,	Lease (leaseNo, paymentMethod, depositPaid, rentStart, rentFinsh, deposit, duration,
privateOwnerNo, buisnessOwnerNo, staffNo, branchNo)	clientNo, propertyNo)
Primary Key propertyNo	Primary Key leaseNo
Foreign Key privateOwnerNo references PrivateOwner(ownerNo)	Alternate Key propertyNo, rentStart
Foreign Key buisnessOwnerNo references BuisnessOwner(buisnessOwnerNo)	Alternate Key clientNo, rentStart
Foreign Key staffNo references Staff(staffNo)	Foreign Key clientNo references Client(clientNo)
Foreign Key branchNo references Branch(branchNo)	Foreign Key propertyNo references PropertyForRent(propertyNo)
	Derived deposit (propertyForRent.rent * 2)
	Derived duration (rentFinish - rentStart)
PrivateOwner (privateOwnerNo, fName, IName, housNo, streetNo, zipCode, city,	BusinessOwner (businessOwnerNo, bName, bType, contactName, housNo, streetNo,
country, telNo)	zipCode, city, country, telNo)
Primary Key privateOwnerNo	Primary Key businessOwnerNo
Alternate Key telNo	Alternate Key bName
	Alternate Key telNo
Advert (propertyNo, newspaperName, advertDate)	Newspaper (newspaperName)
Primary Key propertyNo, newspaperName, advertDate	Primary Key newspaperName
Foreign Key propertyNo references PropertyForRent(propertyNo)	
Foreign Key newspaperName references Newspaper(newspaperName)	
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### **Step 2.2**

- Step 2.2 Validate relations using normalization
  - To validate relations in logical data model using normalization
  - Typically already in 3NF at this point

- To check integrity constraints represented in logical data model
  - Keep database accurate, consistent, complete
- Identify:
  - Required data
  - Unique data
  - Attribute domain constraints
  - Entity integrity
  - Referential integrity

Required data

-Non null attributes contain valid value

Unique data

-Some attributes values are unique

Attribute domain constraints

 Every attribute has domain

- Entity integrity
  - Primary key not null
- Referential integrity
  - Foreign key must reference existing value in parent relation
  - Foreign key null if participation optional
  - Existence constraints

### • Existence constraints

### Delete tuple from child relation

• Referential integrity is unaffected

### Insert tuple into child relation

- Foreign key value will be either set to null or to a value from the primary key attribute of parent relation
- Automatically maintained by foreign key constraint
- Update foreign key of child tuple
  - Foreign key value will be either set to null or to a value from the primary key attribute of parent relation
  - Automatically maintained by foreign key constraint

### • Existence constraints

### Delete tuple from parent relation

- Referential integrity is lost, if there exists a tuple in child relation referencing the deleted tuple of parent relation
- Several strategies:
  - ON DELETE NO ACTION
    - » Prevents a deletion from the parent relation
    - » Safe option
  - ON DELETE CASCADE
    - » Automatically delete all related tuples in child relation
    - » Dangerous option
  - ON DELETE SET NULL
    - » Sets to null the related foreign key values in child relation
  - ON DELETE SET DEFAULT
    - » Sets to default the related foreign key values in child relation

- Existence constraints
  - Insert tuple into parent relation
    - Referential integrity is unaffected
  - Update primary key of parent tuple
    - Referential integrity is lost, if there exists a tuple in child relation referencing the updated tuple of parent relation
    - Strategy:
      - ON UPDATE CASCADE (Typical option)

### Document all integrity constraints

 Document all integrity constraints in the data dictionary for consideration during physical design.

### **Existence constraints for relations of** *DreamHome*

Staff ( staffNo, fName, IName, position, gender, DOB, salary, managerNo, mgrStartDate )
Primary Key staffNo
Foreign Key managerNo references Staff(staffNo) ON UPDATE CASCADE ON DELETE SET NULL

**Owner** (ownerNo, fName, IName, telNo ) **Primary Key** ownerNo

PropertyForRent (propertyNo, address, street, city, postcode, type, rooms, rent, ownerNo, staffNo)
 Primary Key propertyNo
 Foreign Key ownerNo references Owner(ownerNo) ON UPDATE CASCADE ON DELETE NO ACTION
 Foreign Key staffNo references Staff(staffNo) ON UPDATE CASCADE ON DELETE SET NULL

Client (clientNo, fName, IName, prefType, maxRent ) Primary Key clientNo Foreign Key staffNo references Staff(staffNo) ON UPDATE CASCADE ON DELETE NO ACTION

ClientTelNo (clientNo, telNo) Primary Key clientNo, telNo Foreign Key clientNo references Client(clientNo)

#### **Existence constraints for relations of** *DreamHome*

Advertise (propertyNo, newspaperName, dateAdvert, cost ) Primary Key propertyNo, newspaperName, dateAdvert Foreign Key propertyNo references PropertyForRent(propertyNo) ON UPDATE CASCADE ON DELETE NO ACTION Foreign Key newspaperName references Newspaper (newspaperName) ON UPDATE CASCADE ON DELETE NO ACTION Newspaper ( newspaperName, address, telNo, contactName ) Primary Key newspaperName Alternate Key telNo

**SignLease** (leaseNo, signDate, paymentMethod, depositPaid, rentStart, rentFinish, deposit, duration, clientNo, propertyNo)

Primary Key leaseNo, signDate, clientNo, propertyNo

Foreign Key clientNo references Client(clientNo) ON UPDATE CASCADE ON DELETE NO ACTION

**Foreign Key** propertyNo **references** PropertyForRent(propertyNo) ON UPDATE CASCADE ON DELETE NO ACTION

- **Derived** deposit (PropertyForRent.rent \* 2)
- **Derived** duration (rentFinish rentStart)

Viewing ( clientNo, propertyNo, viewDate, comment ) Primary Key clientNo, propertyNo Foreign Key clientNo references Client(clientNo) ON UPDATE CASCADE ON DELETE NO ACTION Foreign Key propertyNo references PropertyForRent(propertyNo) ON UPDATE CASCADE ON DELETE NO ACTION 38

#### **Conceptual Data Model of all Users' Views**



#### Build the Logical Data Model (Relation Diagram) of all Users' Views for DreamHome



### **Chapter Summary**

- Logical database design is the process of constructing a model of the data used in an enterprise based on a specific data model but independent of a particular DBMS and other physical considerations
- A **logical data model** includes *relational schema*, and supporting documentation such as the *data dictionary*, which is produced throughout the development of the model